

Formulation of the Learning Problem

① Concept and function learning (realizable case)

\mathcal{X} - feature space

\mathcal{C} - concept class (collection of subsets of \mathcal{X})

a

a

 ... all images of lowercase 'a'

o

u

 ... negative examples

$X_1, \dots, X_n \stackrel{iid}{\sim} P$ (instances)

$P \in \mathcal{P}$ (class of distributions)

$C^* \in \mathcal{C}$ (unknown concept)

training data: $(X_i, Y_i) \quad i=1, \dots, n$

$$Y_i = \mathbb{1}_{\{X_i \in C^*\}}$$

$Z_i = (X_i, Y_i)$ - instance-label pair

learning algo: $Z^n = (Z_1, \dots, Z_n) \xrightarrow{A_n} \hat{C}_n \in \mathcal{C}$

$$\mathcal{A} = \{A_n\}_{n=1}^{\infty}$$

$$A_n: (\mathcal{X} \times \{0, 1\})^n \rightarrow \mathcal{C}$$

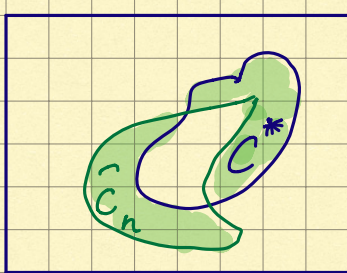
$\hat{C}_n = A_n(Z_1, \dots, Z_n)$ - hypothesis returned by \mathcal{A} on Z^n
- training

$X \sim P$ (indep. of X_1, \dots, X_n)

$$\hat{y} = \mathbb{1}_{\{x \in \hat{C}_n\}}$$

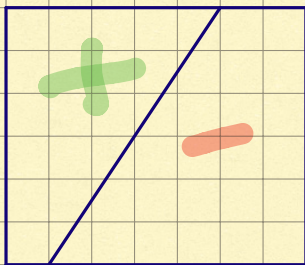
Misclassification: $x \notin C^*, x \in \hat{C}_n$
 $x \in C^*, x \notin \hat{C}_n$

$$\{x \in \mathcal{X} : x \in C^* \cap \hat{C}_n^c \text{ or } x \in (C^*)^c \cap \hat{C}_n\} \\ =: \hat{C}_n \Delta C^* \quad (\text{sym. diff.})$$

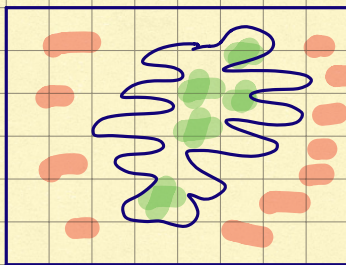


$P[\hat{C}_n \Delta C^*]$ - classification error
↳ r.v. (depends on Z^n)

"Learnability" (\rightarrow complexity (richness of concepts



"simple"



"complex"

Probably Approximately Correct (PAC):

property of a learning algo for a given task

$$\text{data } (Z^n) \xrightarrow{A_n} \hat{C}_n$$

P_C^n - dist. of n iid labeled pts when $C^* = C$
 $P[\hat{C}_n \Delta C^*] \leq \epsilon$ (approx. correct)

$$P_{C^*}^n \left\{ z^n : P[\bar{c}_n \Delta C^*] > \epsilon \right\} < \delta$$

- probably approx. correct

(L. Valiant, D. Angluin 1980s)

Notation : fix $C \in \mathcal{C}$

$$(X, Y) : X \sim P, Y = \mathbb{1}_{\{X \in C\}}$$

$$P_C = \text{dist. of } (X, Y)$$

$$P_C^n = \text{dist. of } (X_1, Y_1), \dots, (X_n, Y_n) \stackrel{\text{i.i.d.}}{\sim} P_C$$

PAC learnability : property of a learning task
(\mathcal{P}, \mathcal{C})

PAC learnable $\Rightarrow \exists$ a PAC learning algo A
for $(\mathcal{P}, \mathcal{C})$

$$\text{Fix } A : r_A(n, \epsilon, P) := \sup_{C \in \mathcal{C}} P_C^n \left\{ P[\bar{c}_n \Delta C] > \epsilon \right\}$$

$$\bar{r}_A(n, \epsilon, \mathcal{P}) := \sup_{P \in \mathcal{P}} r_A(n, \epsilon, P)$$

A is PAC if $\lim_{n \rightarrow \infty} \bar{r}_A(n, \epsilon, \mathcal{P}) = 0$.
(at accuracy ϵ)

To prove PAC learnability: produce PAC learning algo

Examples:

1) Any finite concept class is PAC-learnable

$$\mathcal{C} = \{c_1, c_2, \dots, c_m\} \quad m < \infty$$

Claim: \exists a learning algo A s.t.

$$\bar{r}_A(n, \epsilon, \mathcal{P}) \leq M(1-\epsilon)^n$$

where \mathcal{P} is the class of all distributions on \mathcal{X}

$n \geq \frac{1}{\epsilon} \log \frac{M}{\delta}$ examples suffices to give

$$\bar{r}_A(n, \epsilon, \mathcal{P}) \leq \delta.$$

Proof sketch: data \mathcal{Z}^n

$$\mathcal{F}(\mathcal{Z}^n) := \{m \in [M] : \forall i, y_i = \mathbb{1}_{\{x_i \in c_m\}}\}$$

$$\hat{m}_n := \min \{m : m \in \mathcal{F}(\mathcal{Z}^n)\}$$

$$\hat{C}_n = c_{\hat{m}_n}$$

"bad set": $\mathcal{B} := \{m : P(C_m \Delta C_{m^*}) > \epsilon\}$

$$P_{C^*}^n[\hat{m}_n \in \mathcal{B}] = \sum_{m \in \mathcal{B}} P_{C^*}^n[\hat{m}_n = m]$$

$\hat{m}_n \in \mathcal{F}(\mathcal{Z}^n) : c_{\hat{m}_n}$ consistent w/ data

$\hat{m}_n \notin \mathcal{B} : P[C_m \Delta C_{m^*}] > \epsilon$

$$P_{C^*}^n [\bar{m}_n = m] \leq P[X_i \notin C_m^* \Delta C_m, \forall i] \leq (1 - \epsilon)^n$$

$(m \in B)$

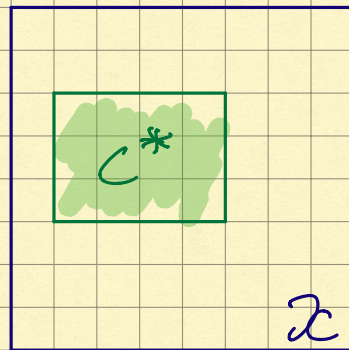
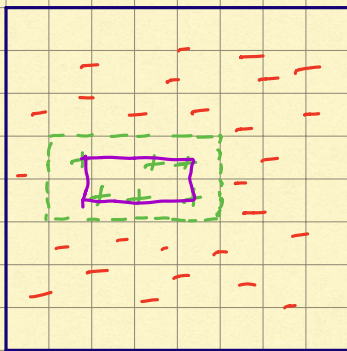
□

2) Axis-aligned rectangles are PAC learnable

$\mathcal{X} = [0, 1]^2$ (unit square)

$\mathcal{C} = \{ [a_1, b_1] \times [a_2, b_2] : a_1, b_1, a_2, b_2 \in [0, 1] \}$

Algo:



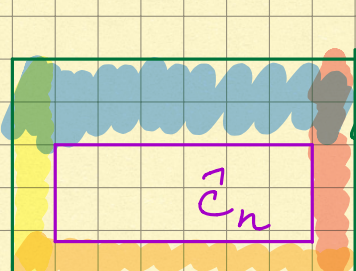
$A_n: \text{data} \rightarrow \hat{C}_n$
 smallest $C \in \mathcal{C}$ that contains all positive examples

Claim: $\bar{r}_A(n, \epsilon, \mathcal{P}) \leq 4(1 - \epsilon/4)^n$

↓
 all dist. on $[0, 1]^2$

(Blumer, Ehrenfeucht, Haussler, Warmuth, 1980s)

idea:



$C^* \Delta \hat{C}_n =$
 ↓ ↓ ↓ ↓
 random sets

show that, w.p. $\geq 1 - 4(1 - \varepsilon/4)^n$,

each of these random rectangles has $P(\cdot) \leq \varepsilon/4$

$$\bar{r}_A(n, \varepsilon, P) \leq 4(1 - \varepsilon/4)^n \leq 4e^{-n\varepsilon/4} \leq \delta$$

$$\text{if } n \geq \frac{4}{\varepsilon} \log \frac{4}{\delta}$$